

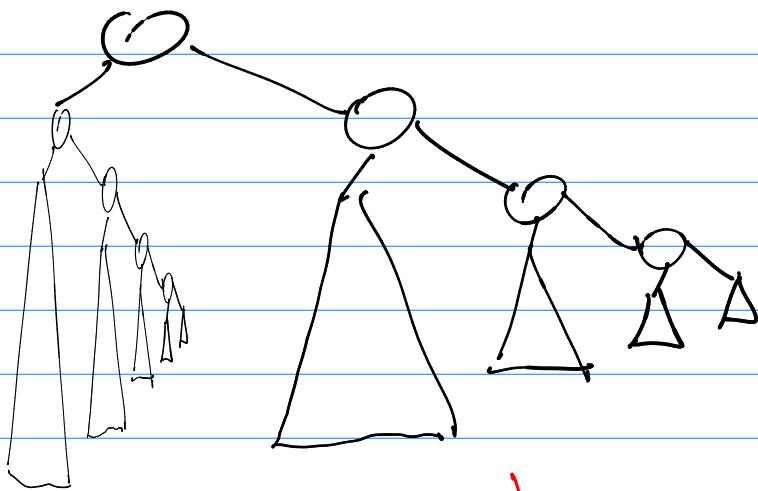
CC3001

# Propiedades

de árboles de Búsqueda

① Altura de AVL  $\in O(\log n)$

Cual es la altura  $h$  de un árbol "AVL" de  $n$  elementos / nodos internos?

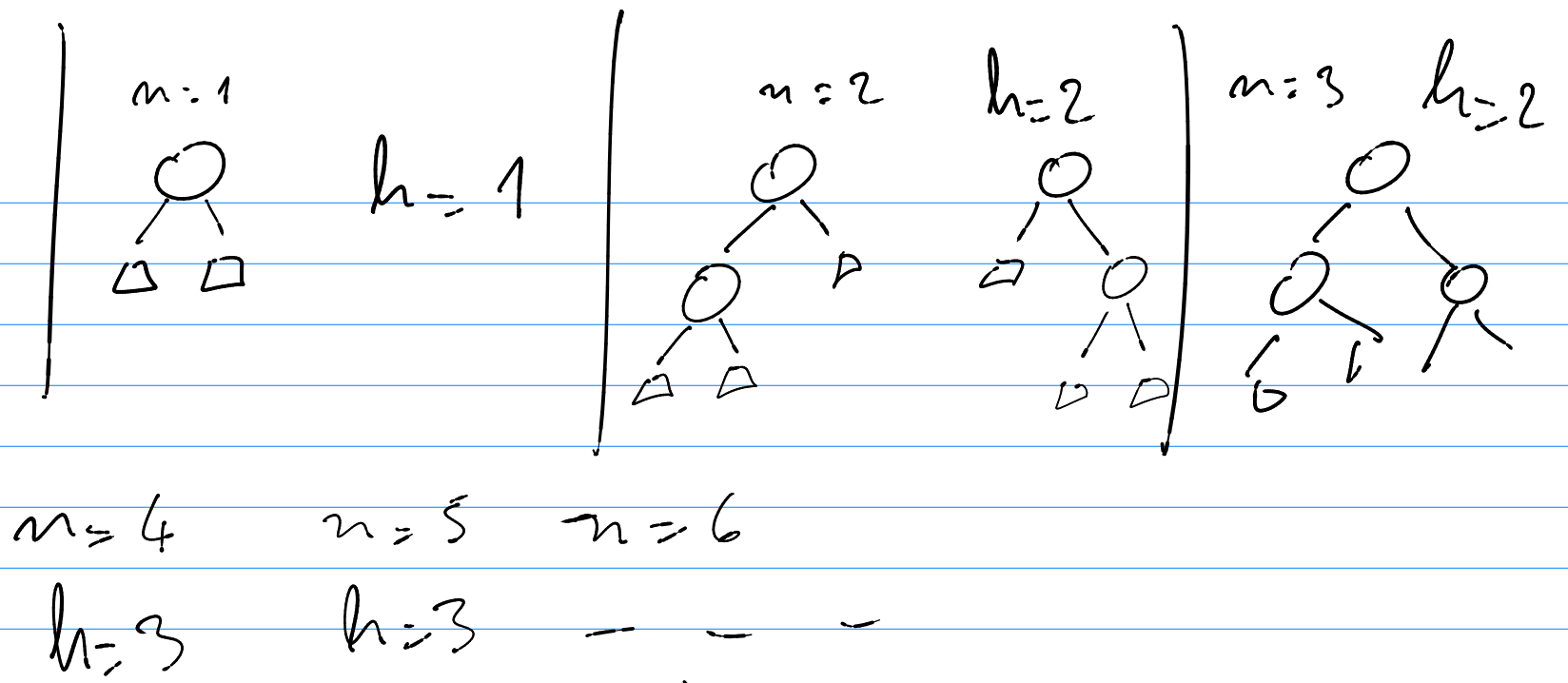


$$h \in O(\log n)$$

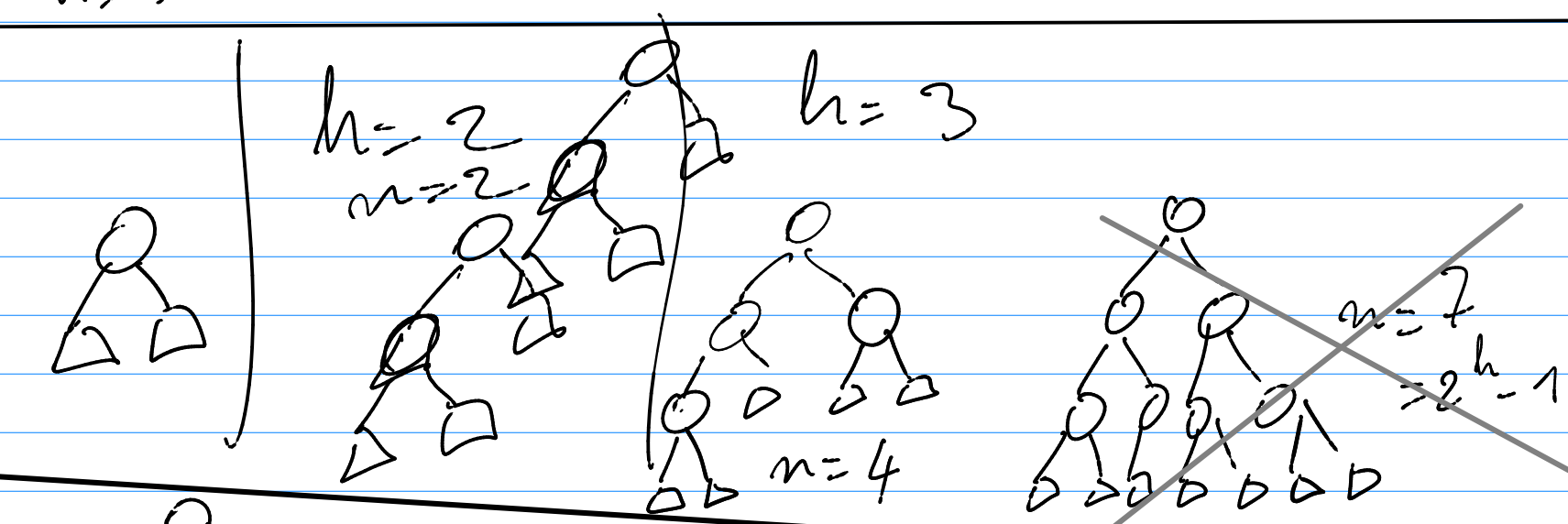
$\equiv$  Cuantos  $n$  nodos tiene un árbol AVL de altura  $h$

$$\Leftrightarrow \begin{cases} \log_2 n < h \leq \log_c n = \frac{\log_2 n}{\log_2 c} \\ c^h = \binom{n}{h} \leq n \leq 2^h - 1 \end{cases}$$

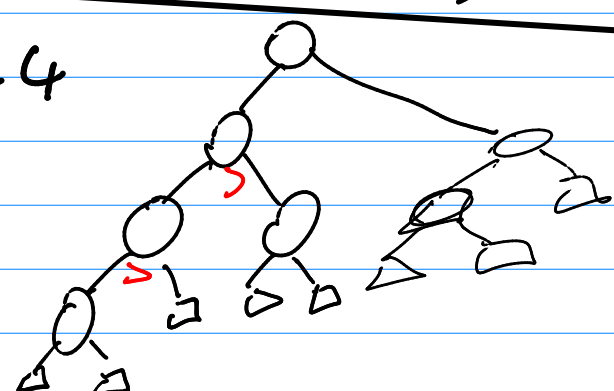
Inducción sobre  $n$ :  
 demuestro "lento"



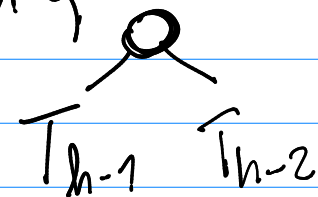
Inducción sobre  $h$   
 $h=1$   
 $n=1$



$h=4$   
 $n=7$



$h$   $T_h$  que minimiza  $n_h$

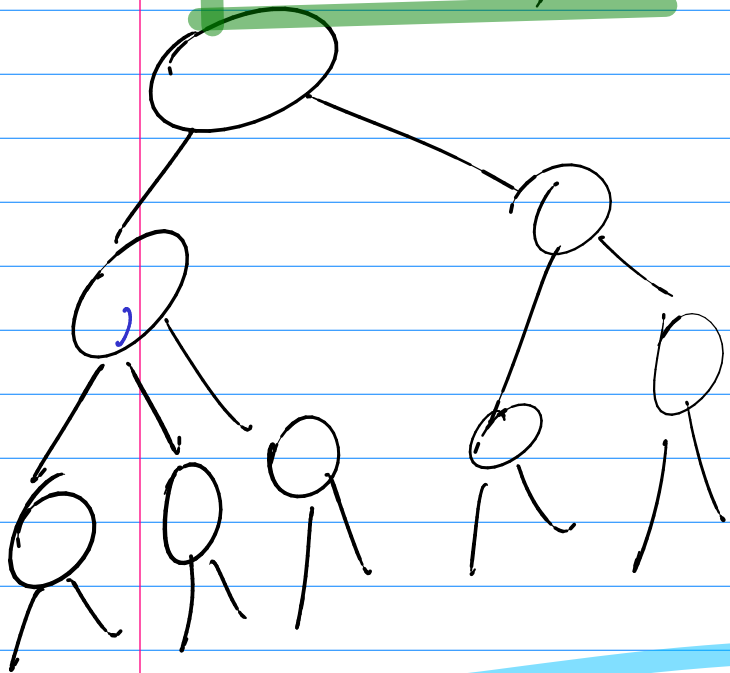


$n_h = 1 + n_{h-1} + n_{h-2}$   
 $n_1 = 1$   
 $n_2 = 2$   
**FIBONACCI**



## ② Extensiones

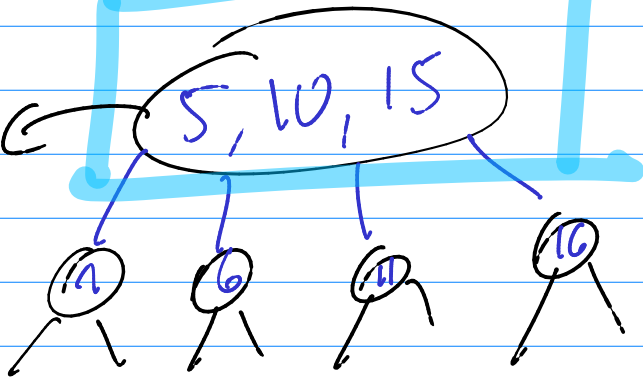
Arbol 2,3)



B-arboles

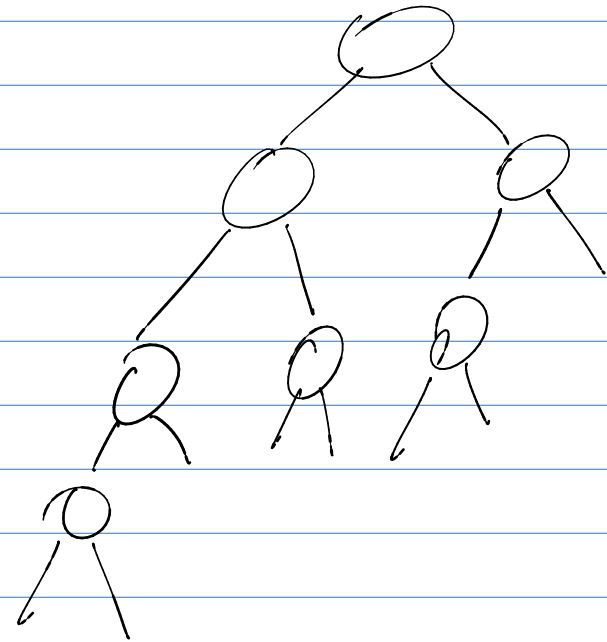
"pagina" de B elementos

$\left[ \frac{B}{2} \dots B \right]$



$\log_B n$

Arboles AVL

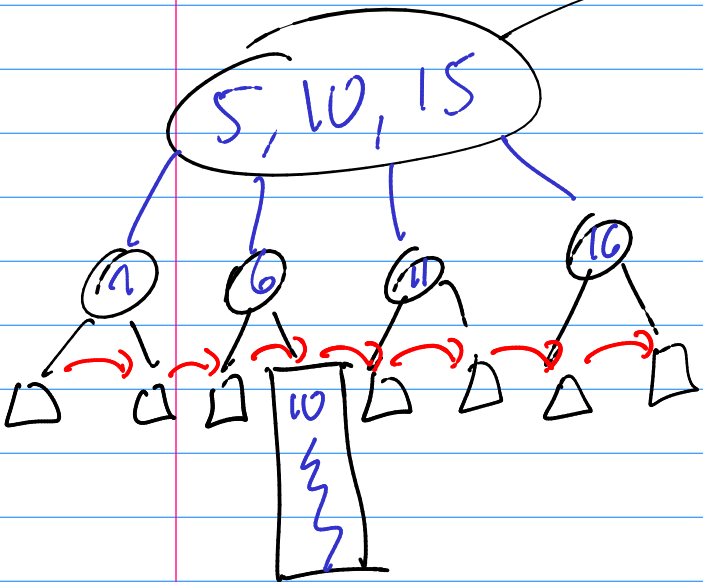


(Podríamos relajarse más con el desequilibrio)

# Arboles B\*

↳ Lista enlazada de hojas

250

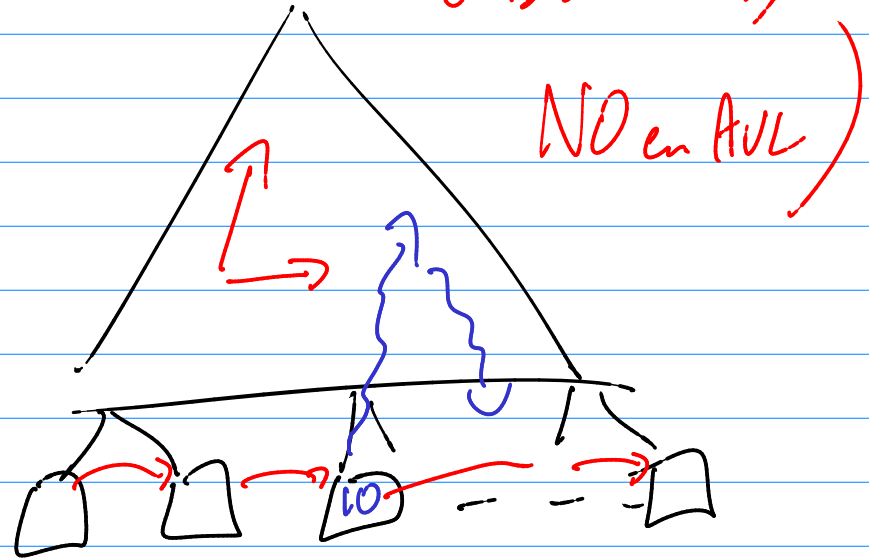


y

# Finger Search Tree

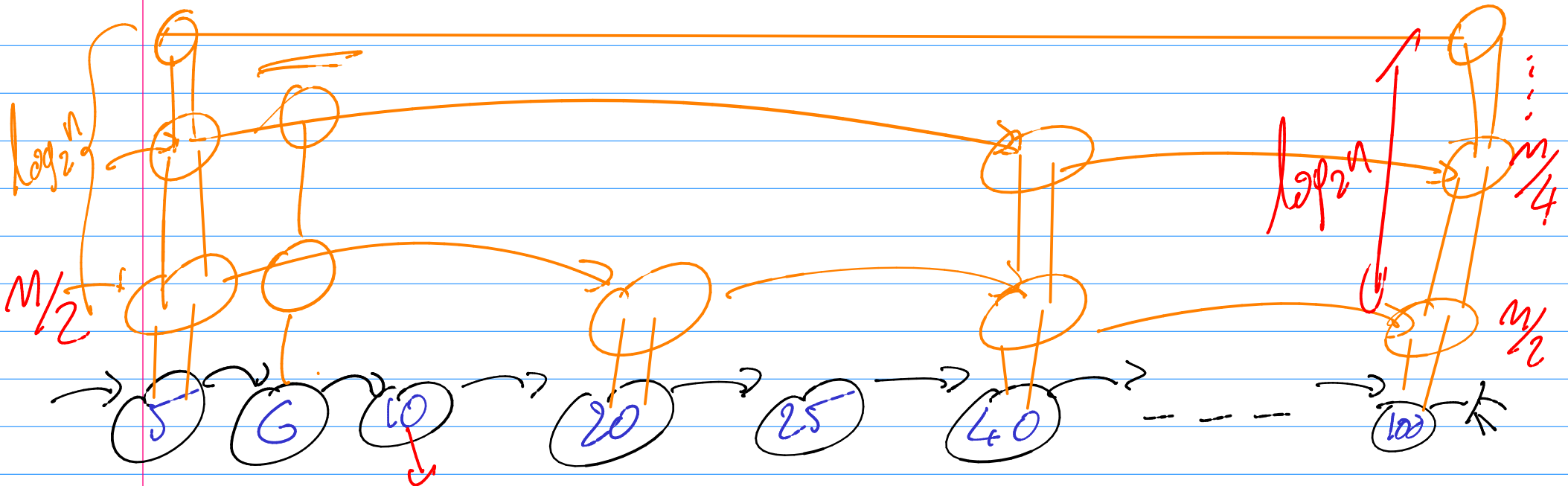
(en Arboles 23 o B arboles,  
o B arboles,

NO en AVL)

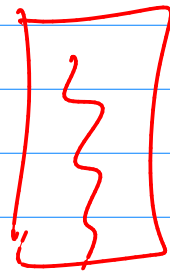


- una lista enlazada
- un arbol indexando la lista

# ③ Skip Lists



A = +1  
B = Pare



Busqueda  $O(\log n)$   
 Insercion \_\_\_\_\_  
 Delecion \_\_\_\_\_

} En Promedio!